A Model of Independence and Overlap for Transactions on Database Schemata

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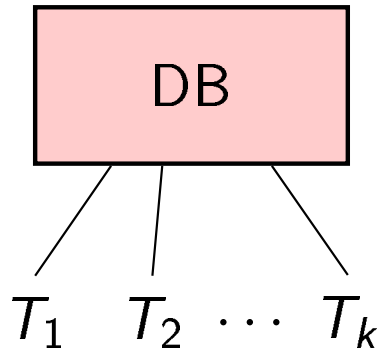

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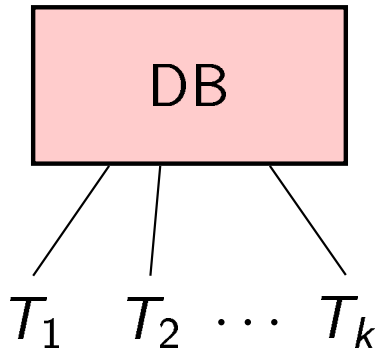

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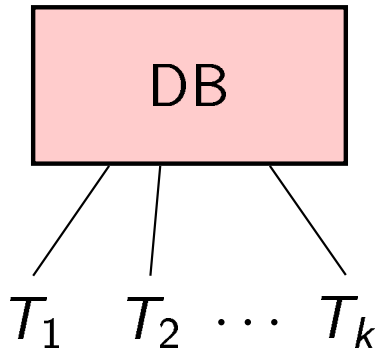

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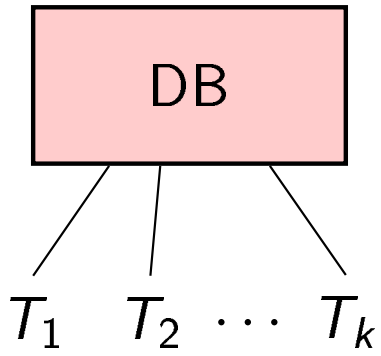

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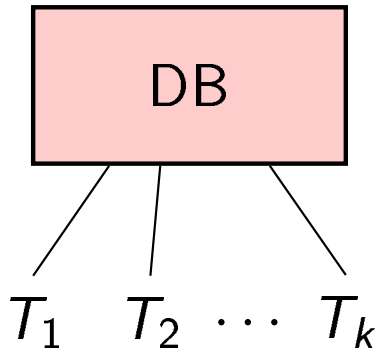

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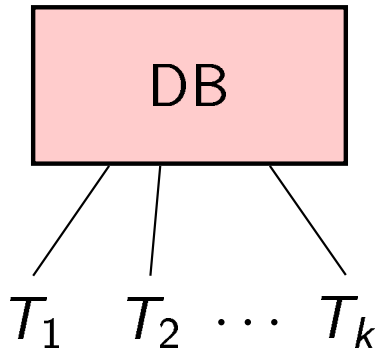

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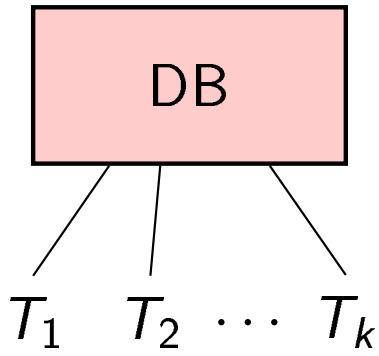

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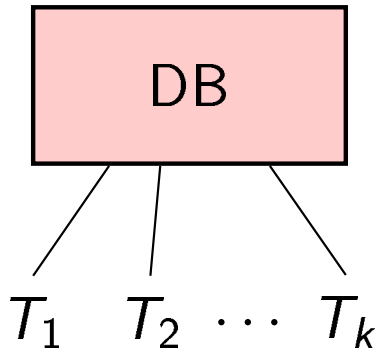

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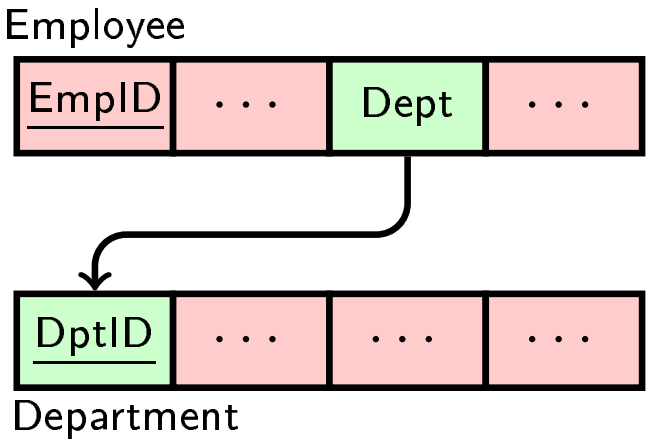

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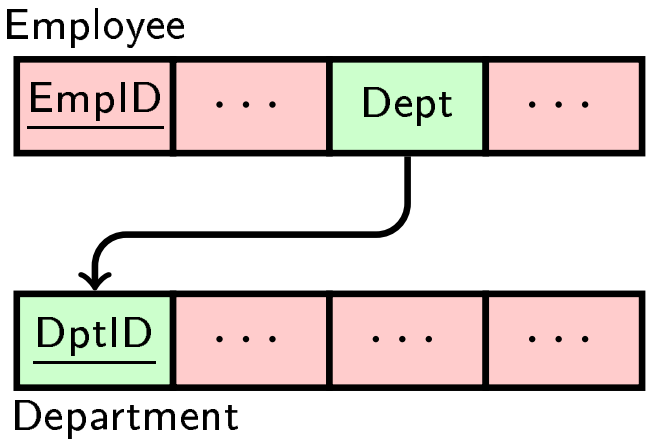

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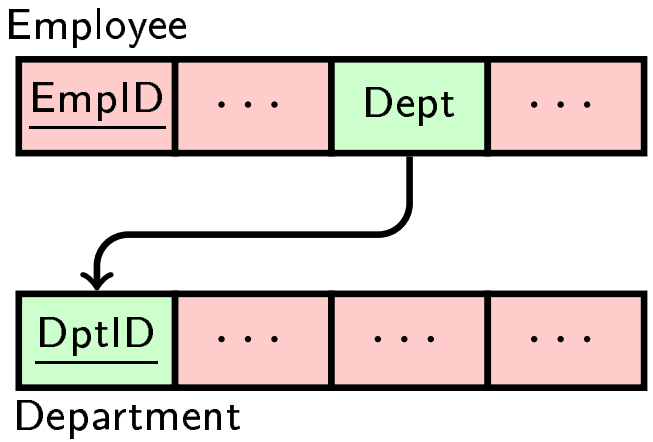

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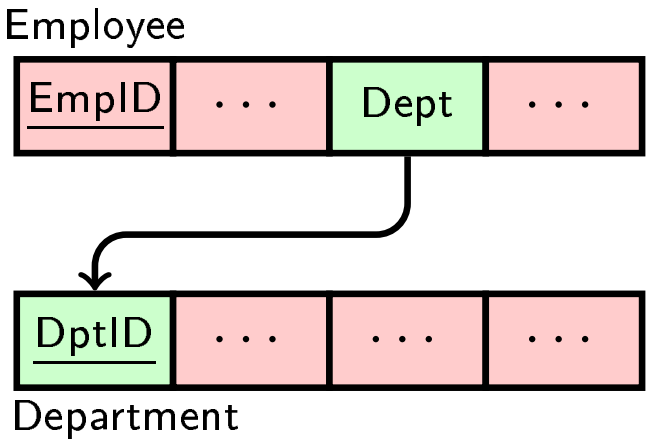

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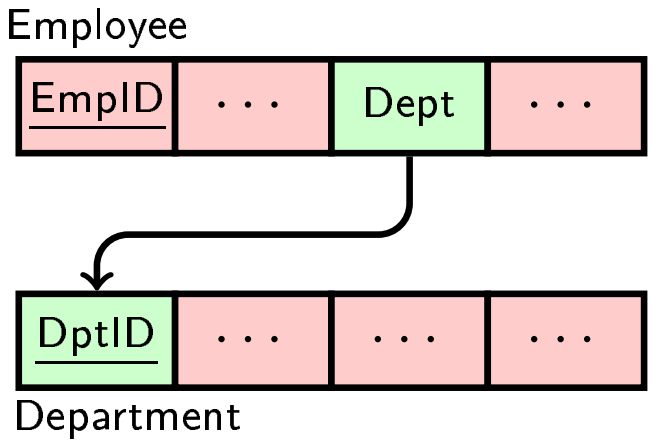

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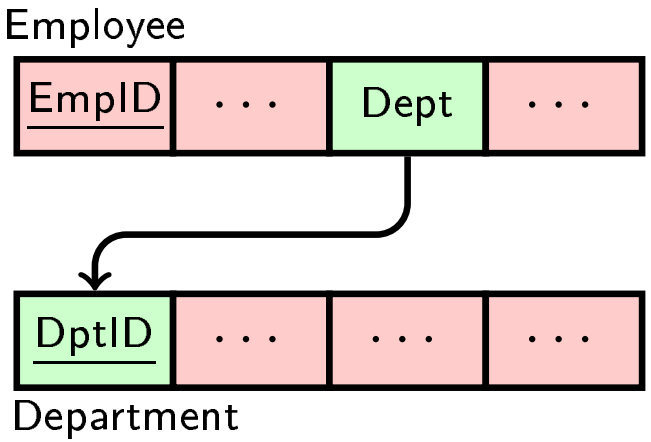

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- Their overlap is only of a very limited read-only nature.

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Focus of this research:

- A fine-grained model of interdependence for data objects.

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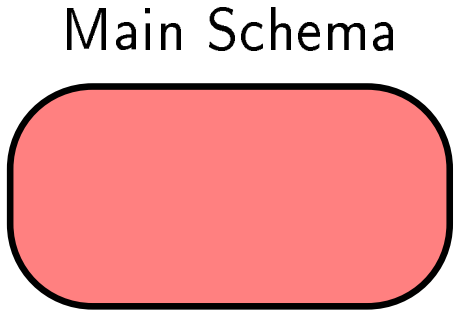

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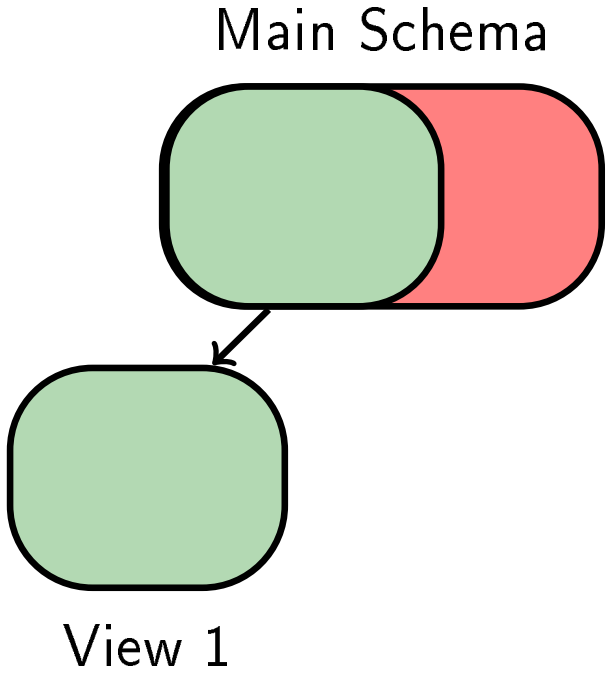

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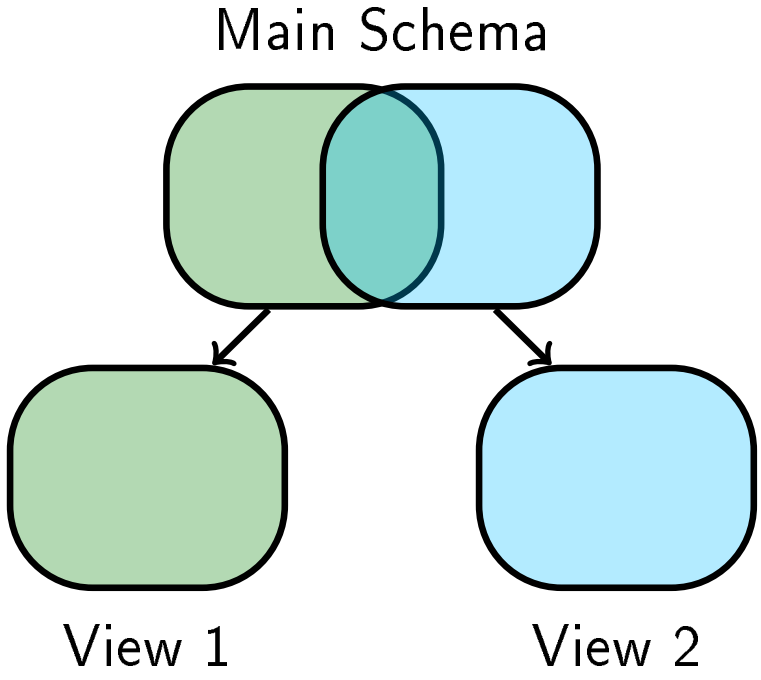

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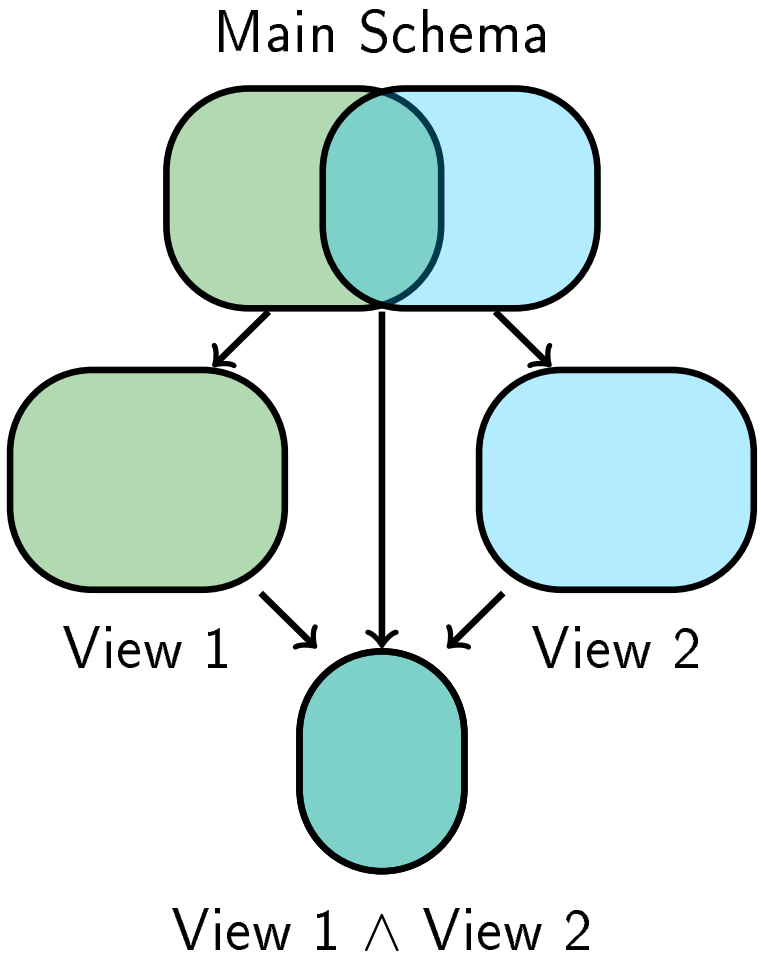

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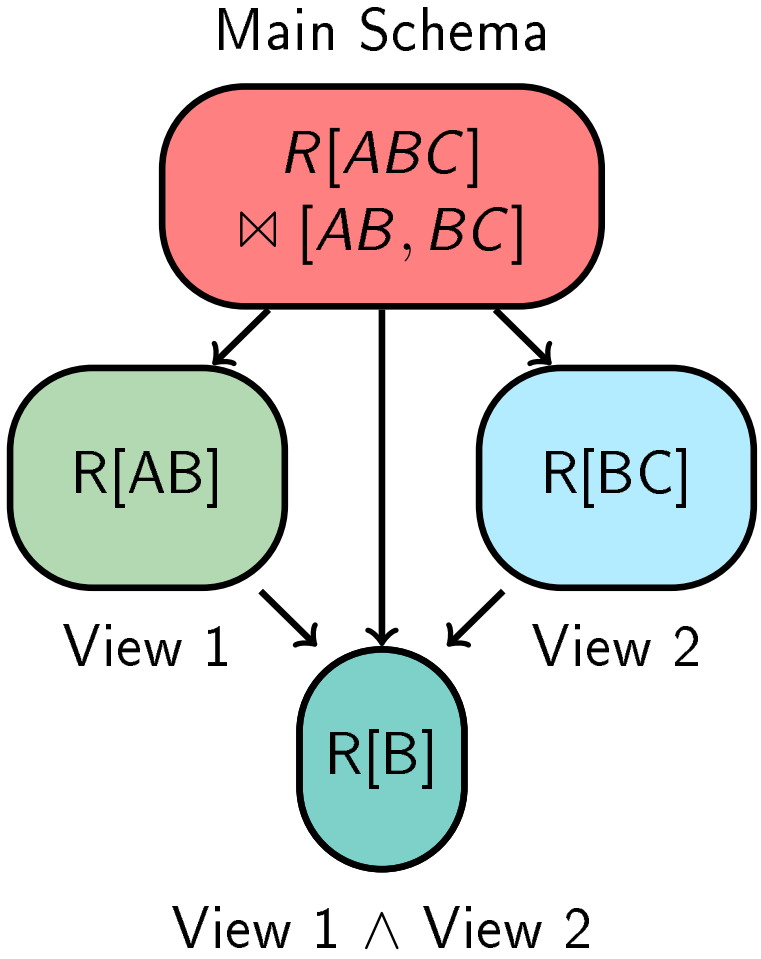

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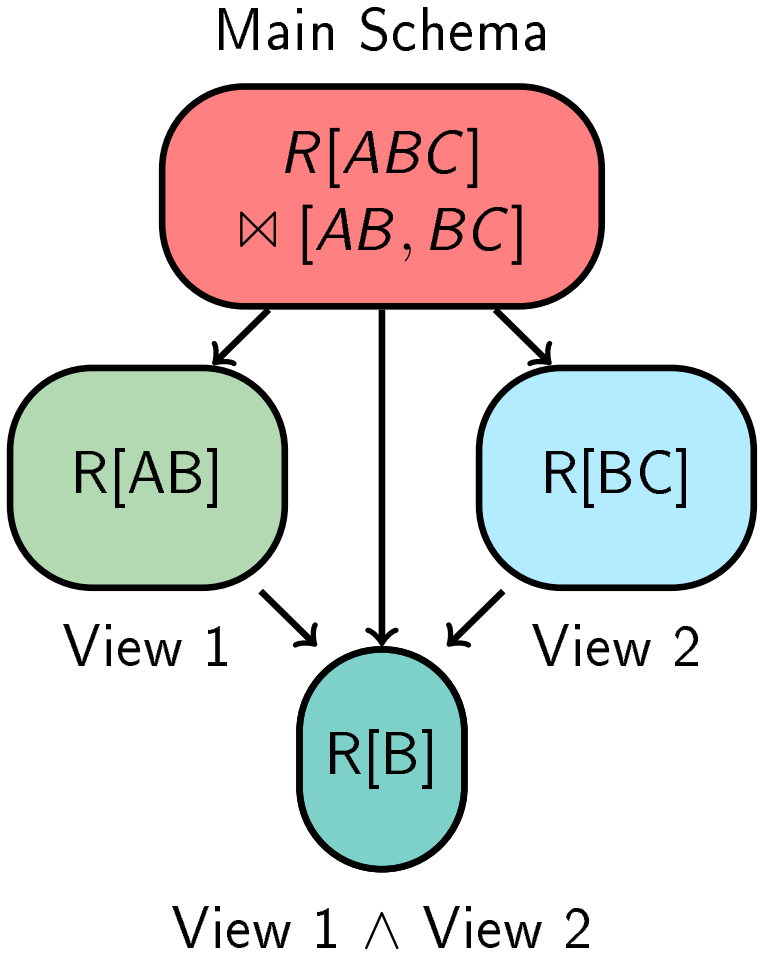

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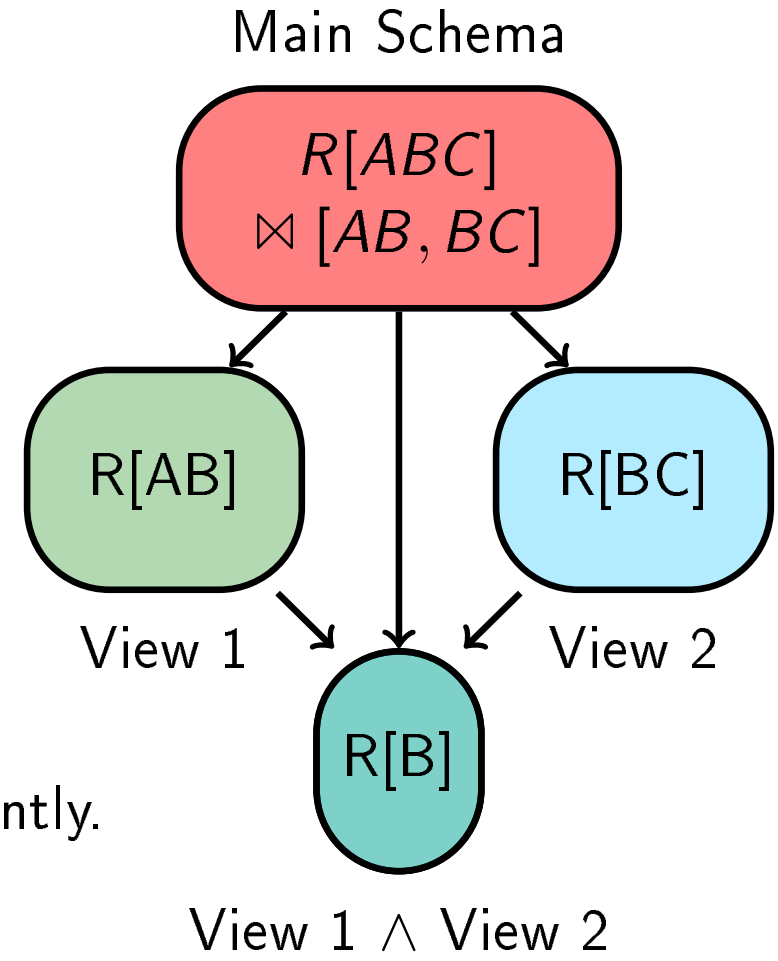

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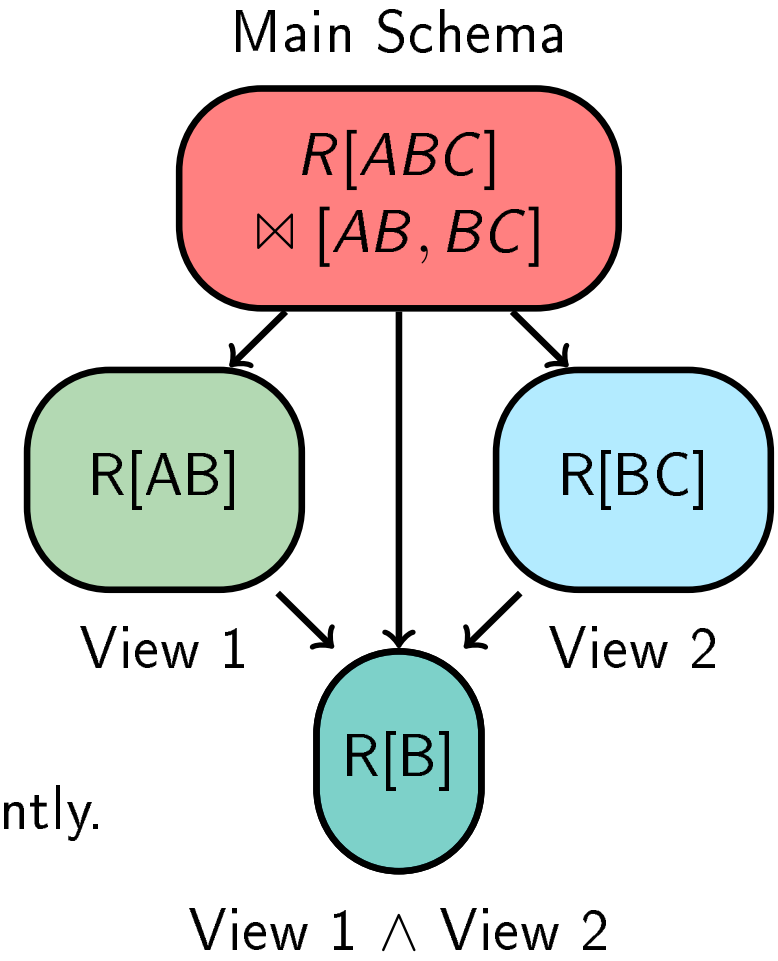

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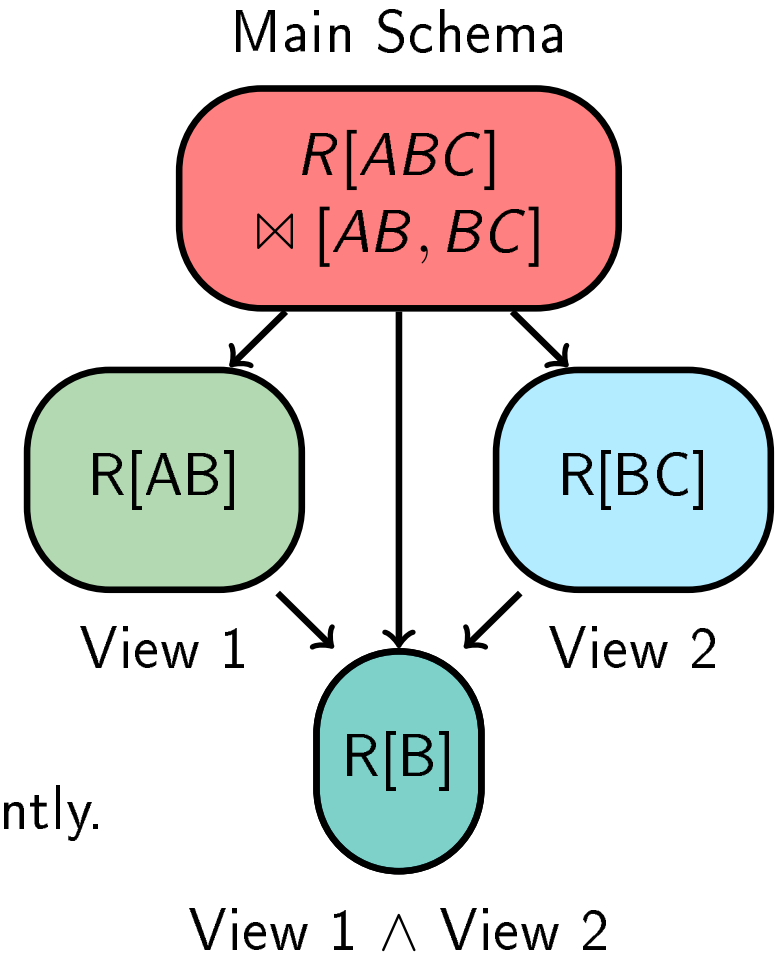

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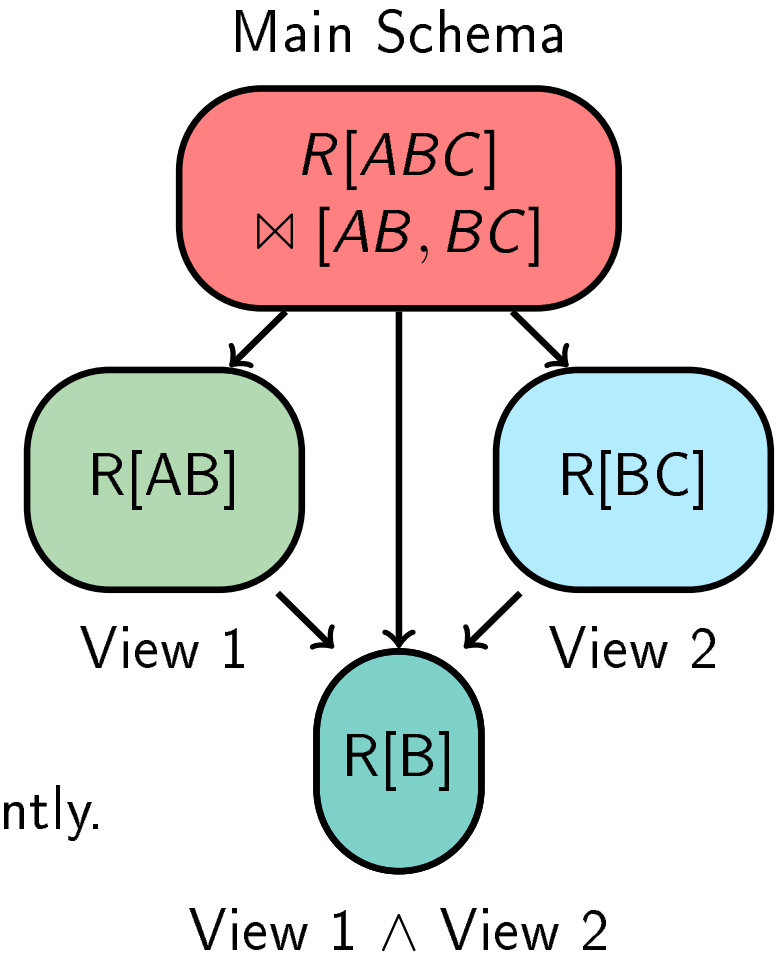

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- This forms the basic idea for independent data objects.

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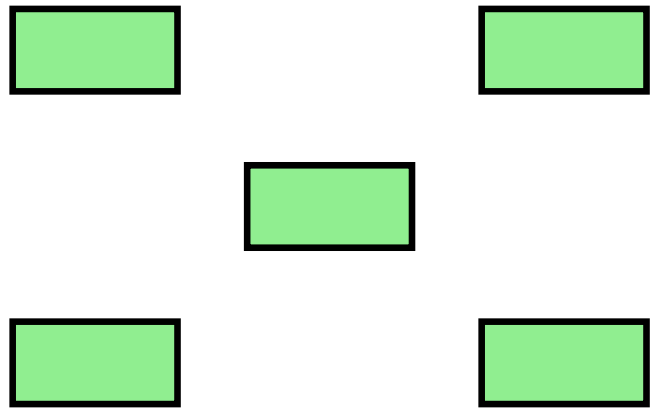

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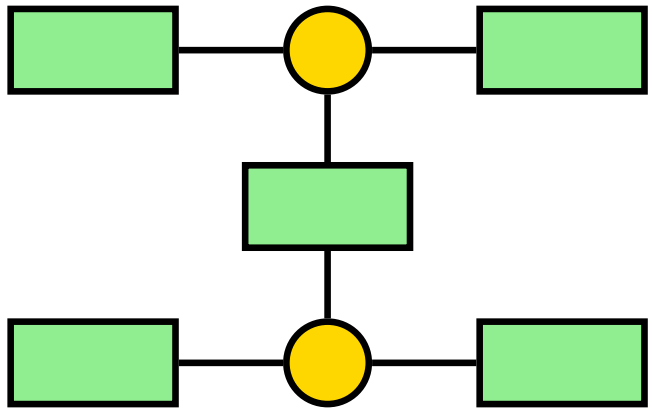

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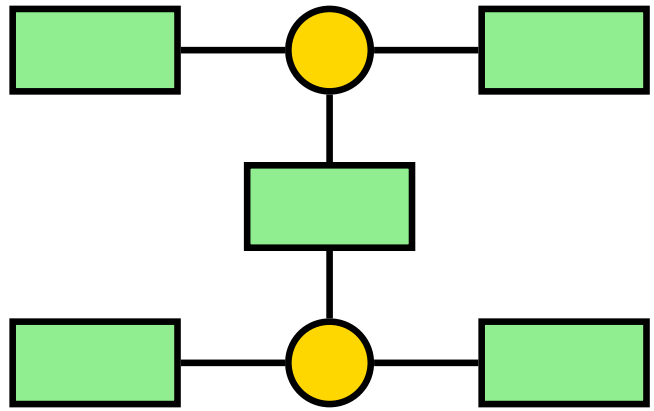

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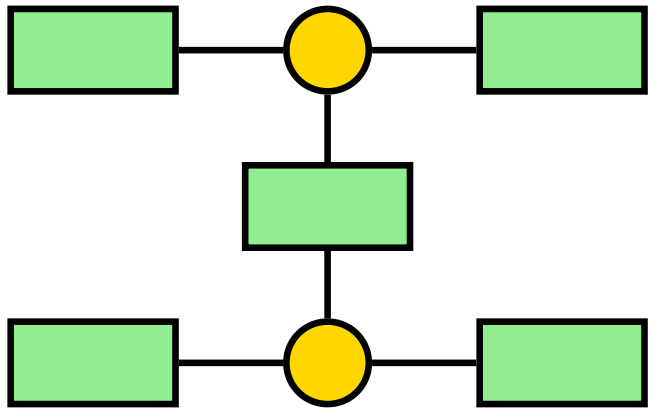

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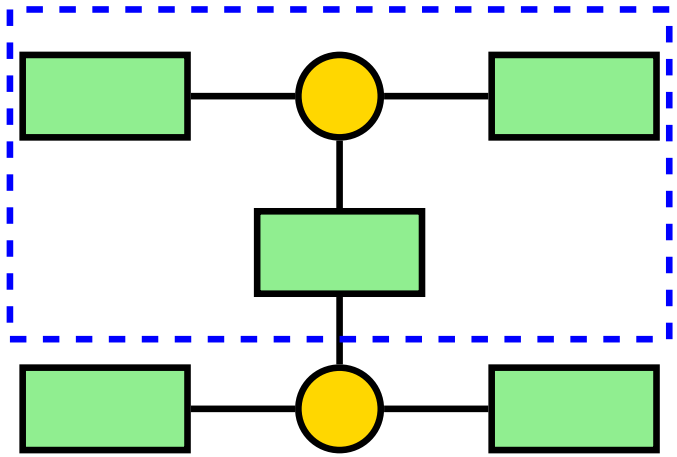

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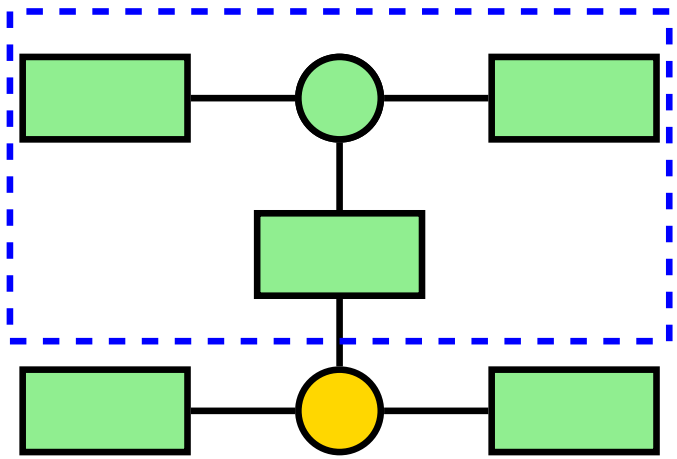

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- To obtain a write claim on a port, all basic components which share that port must be combined into a larger complex object.

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- The remainder of the talk will sketch how these goals are realized.

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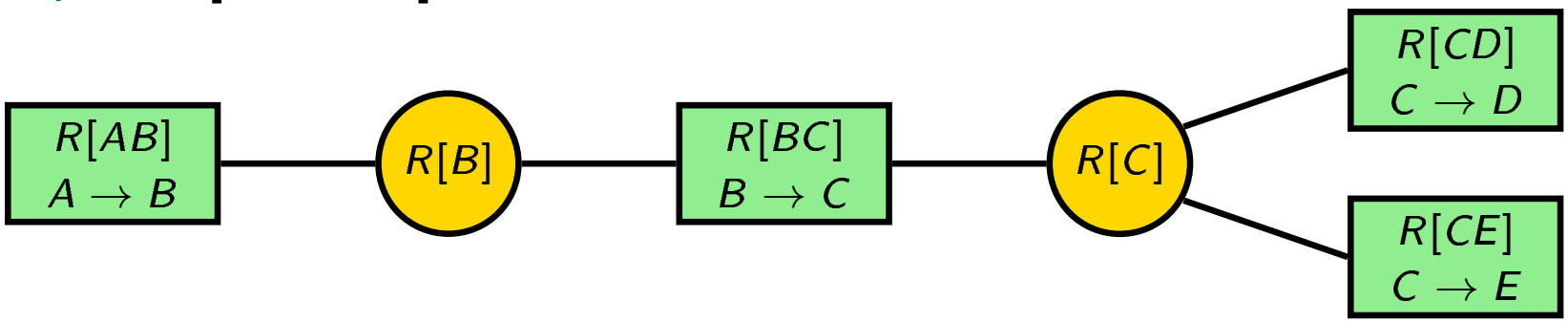

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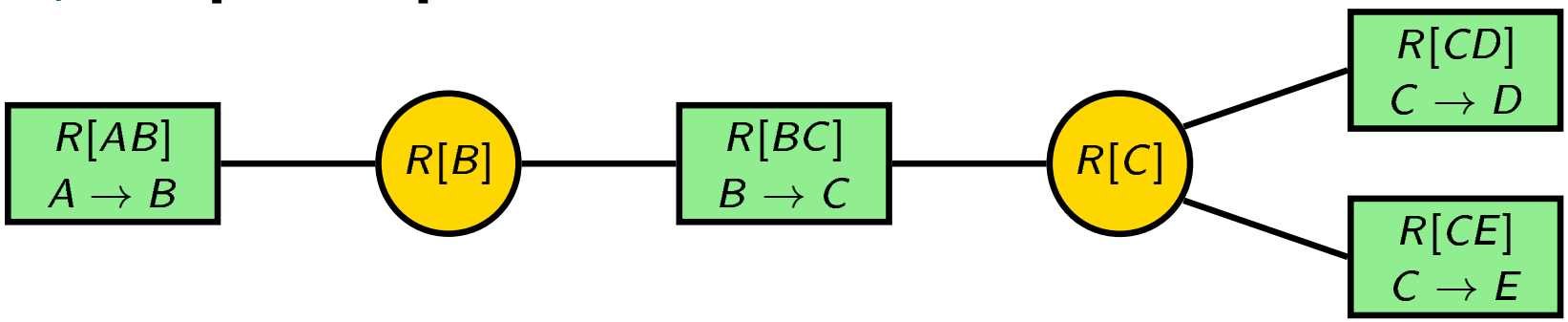

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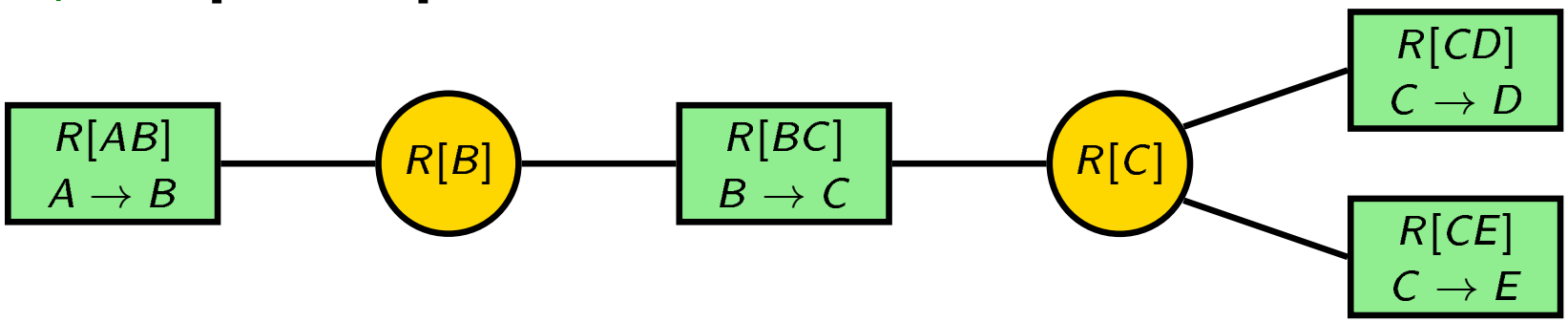

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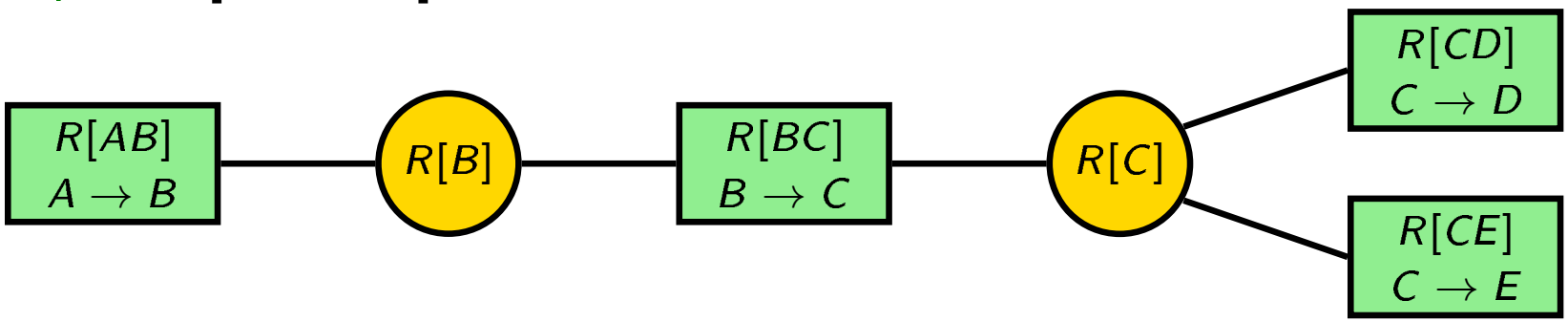

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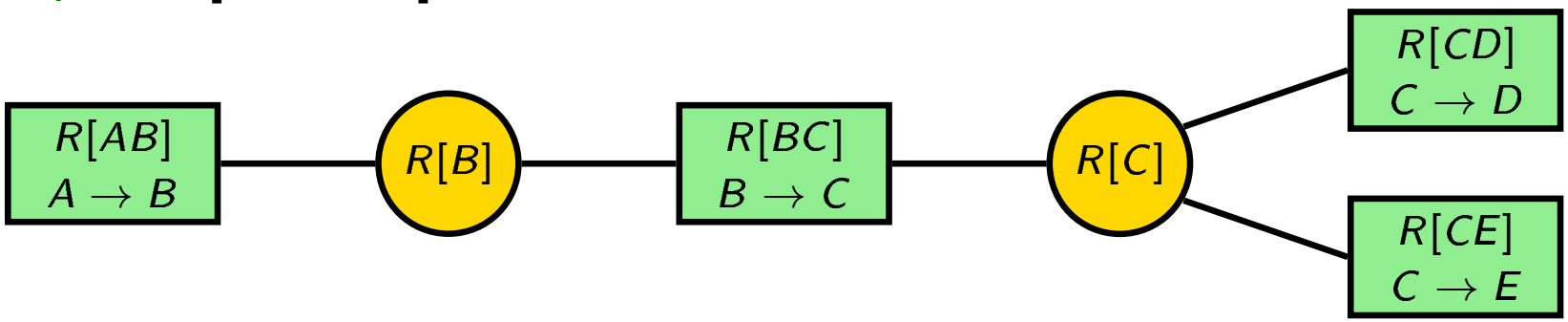

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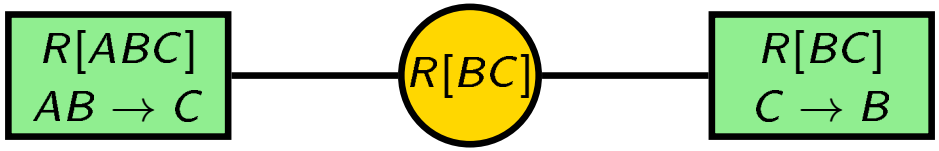

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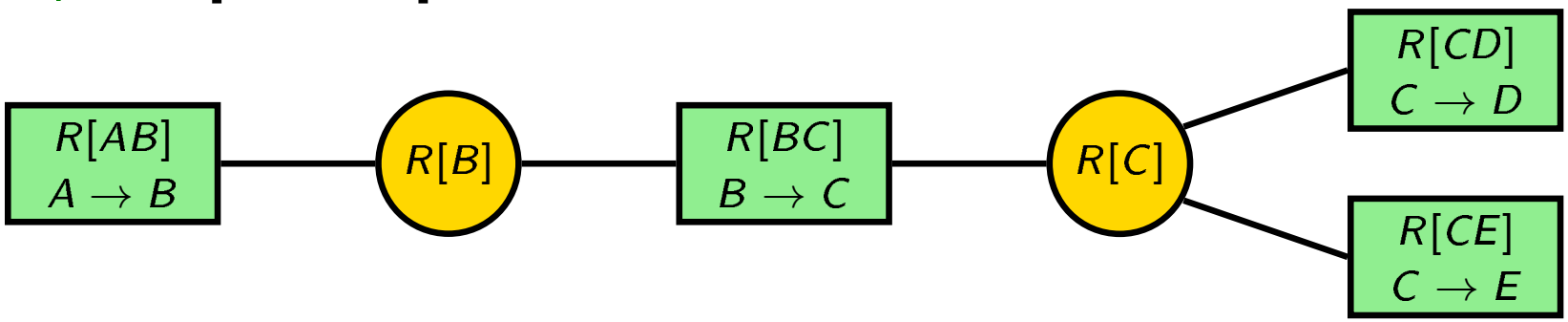

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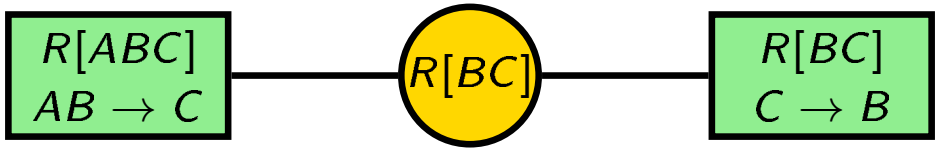

- These are object definitions, not materialized views!!

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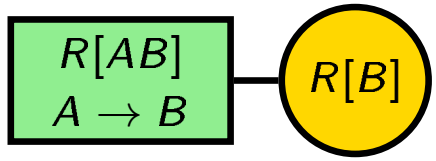

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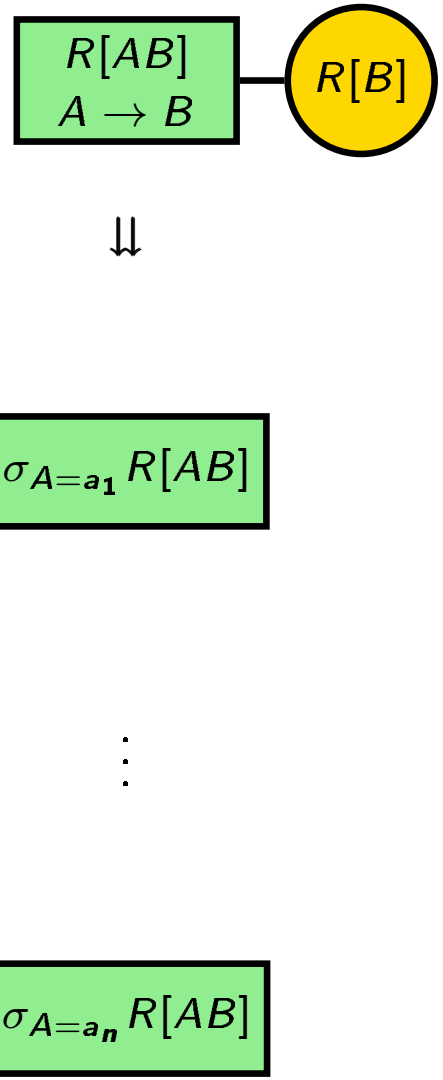

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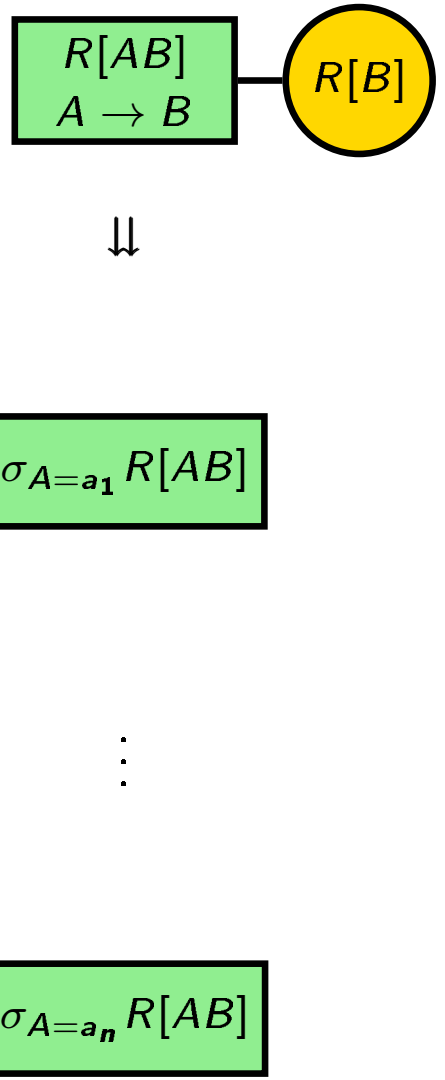

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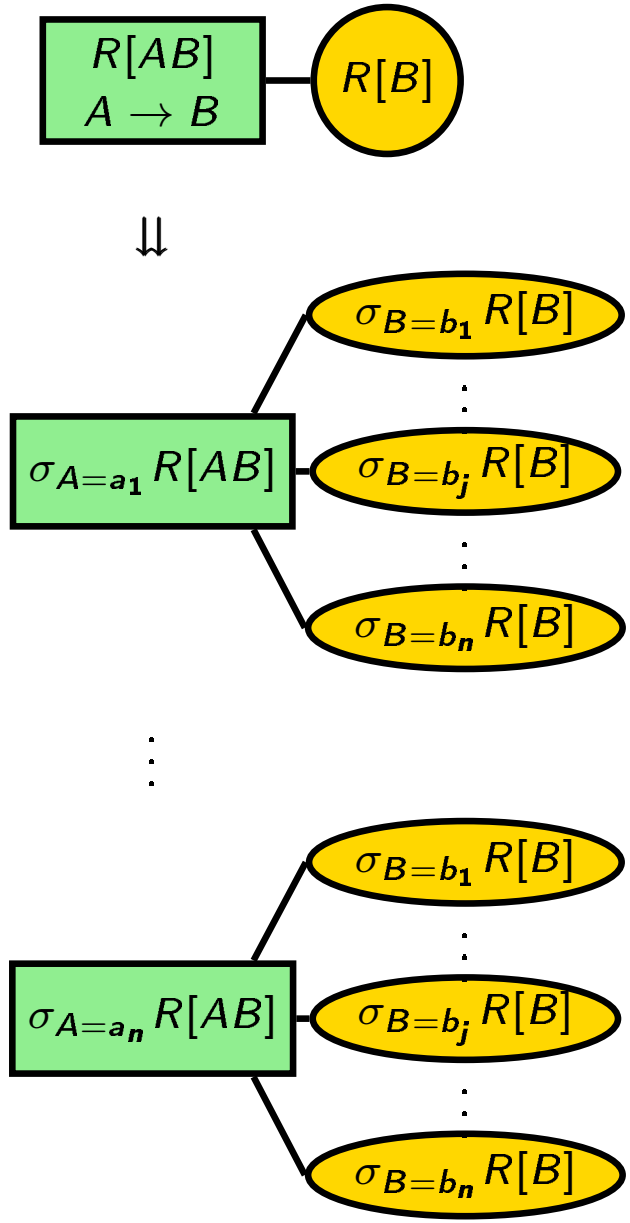

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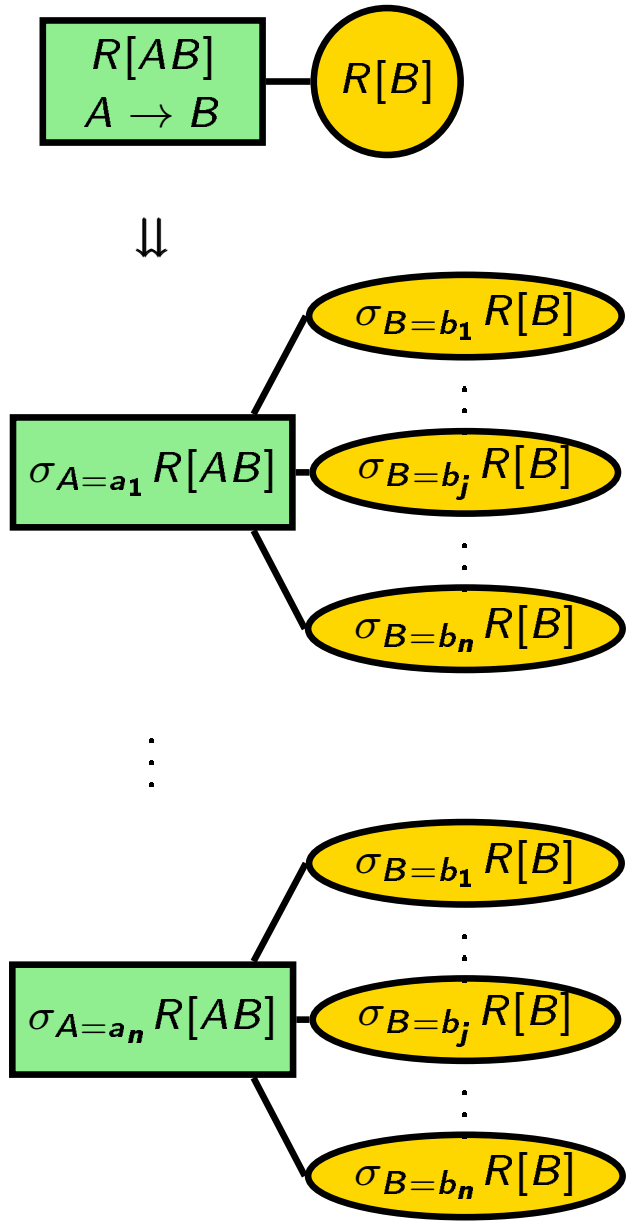

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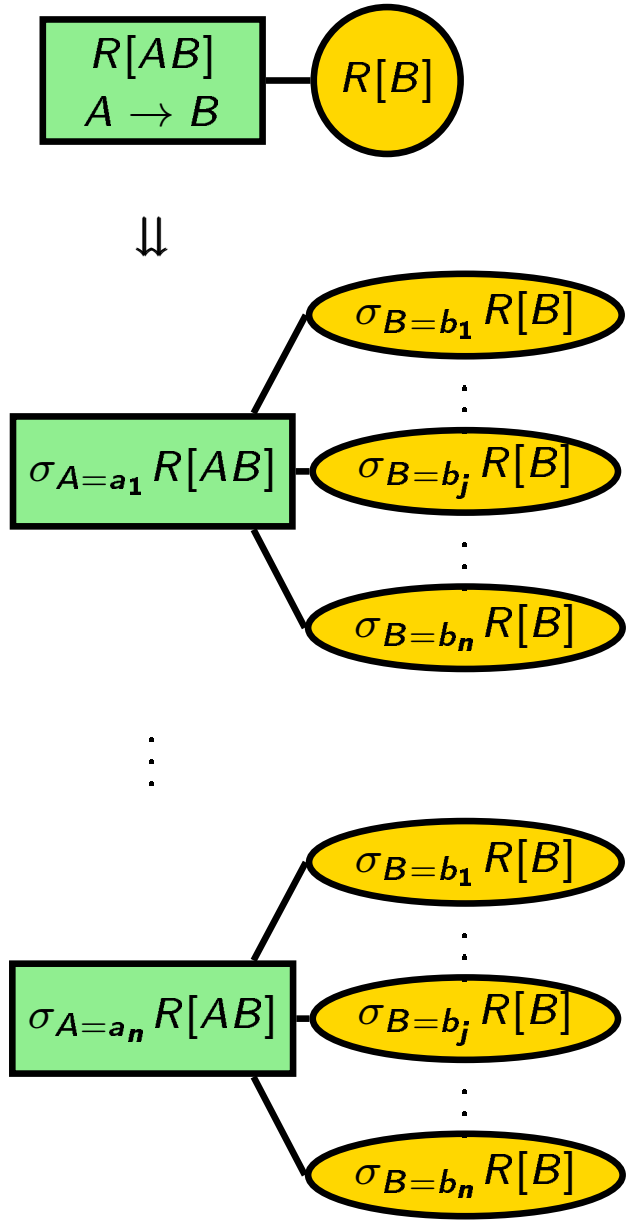

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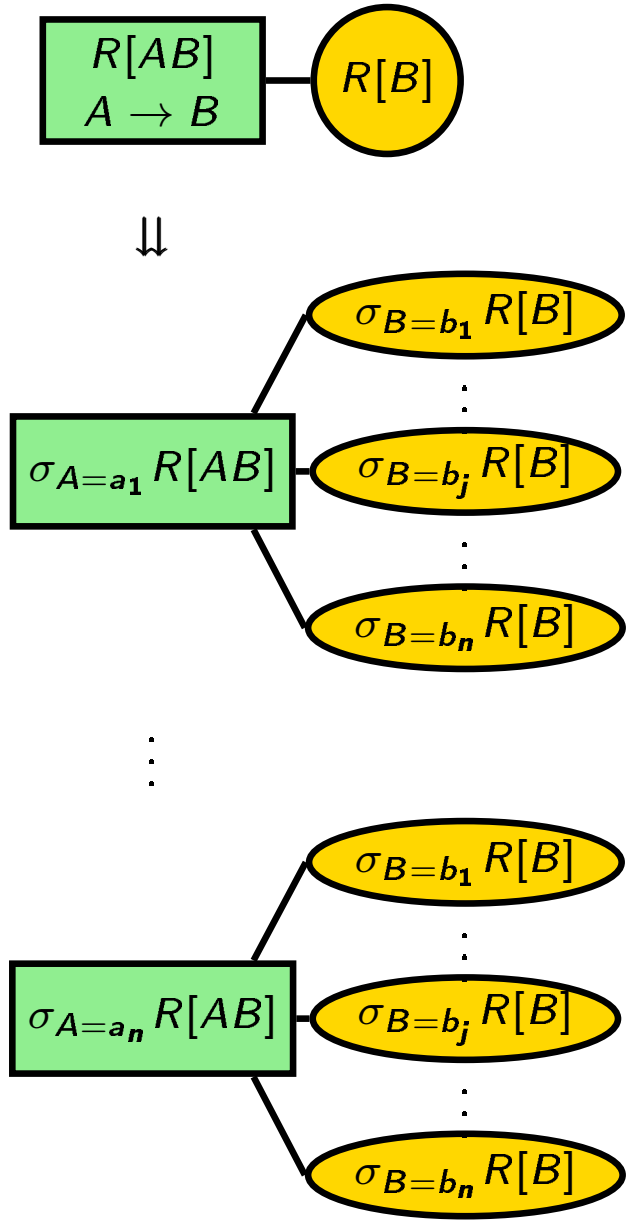

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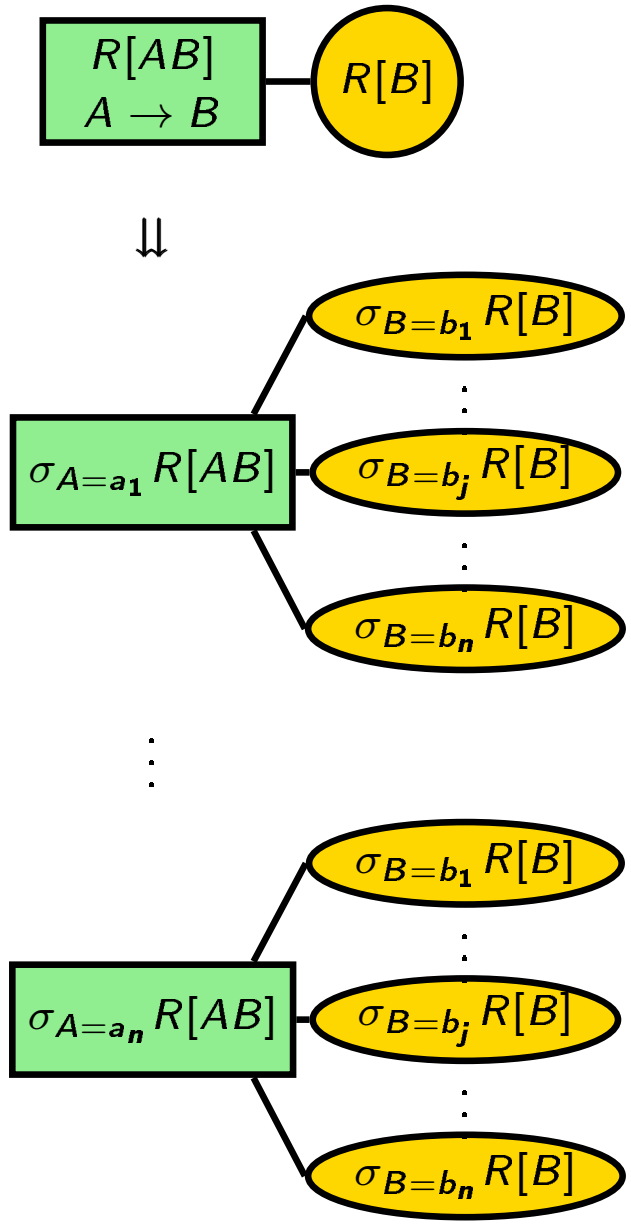

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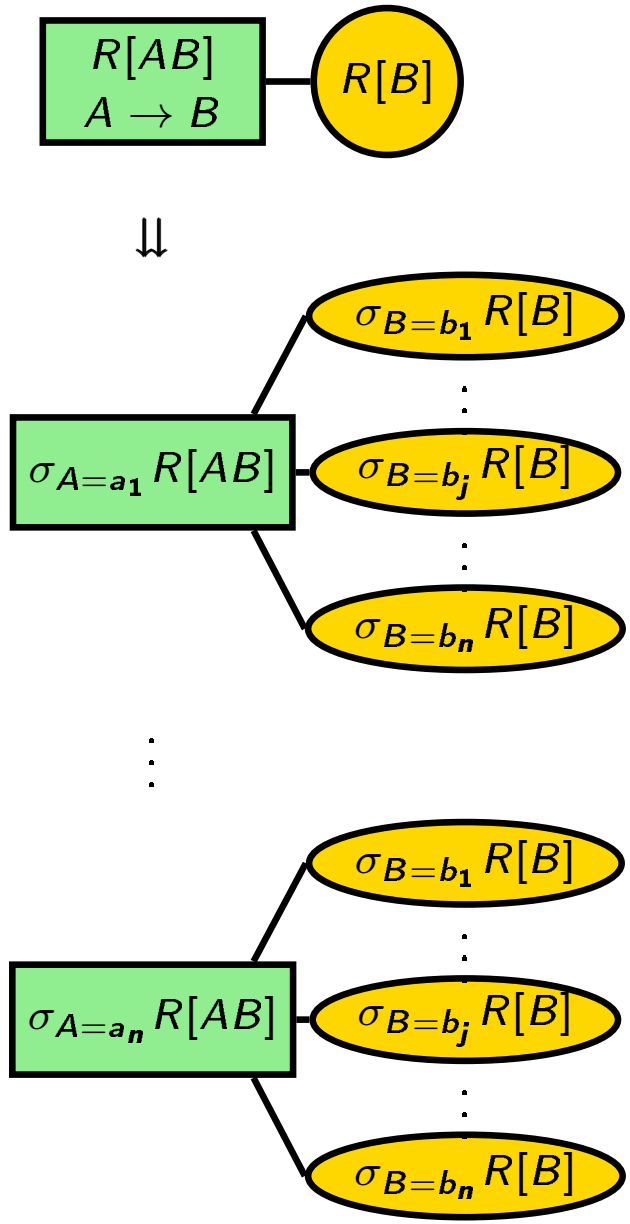

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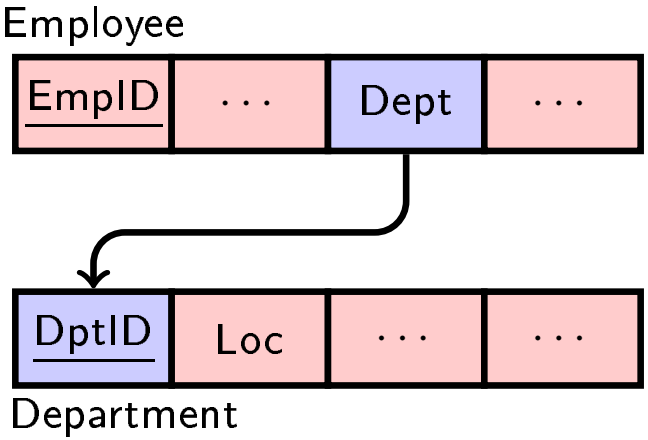

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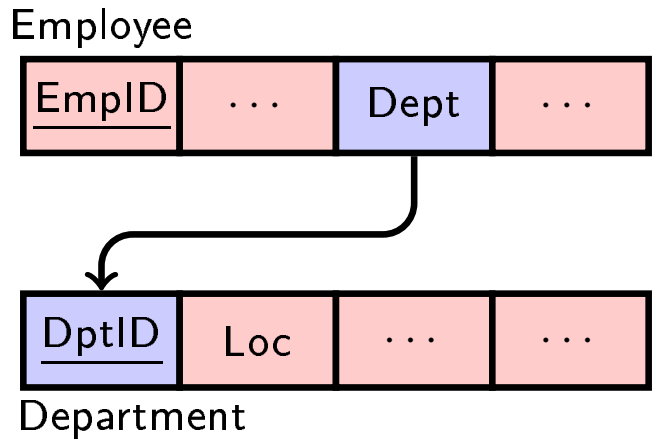

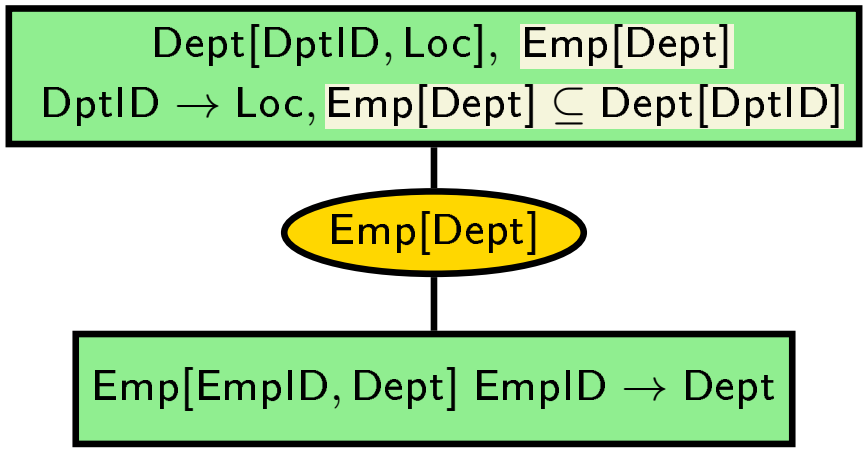

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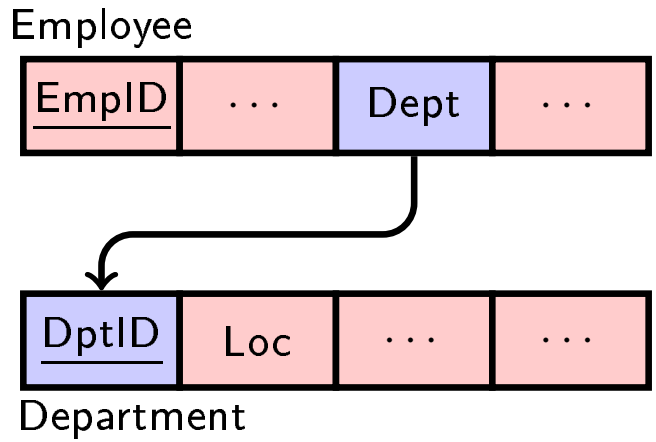

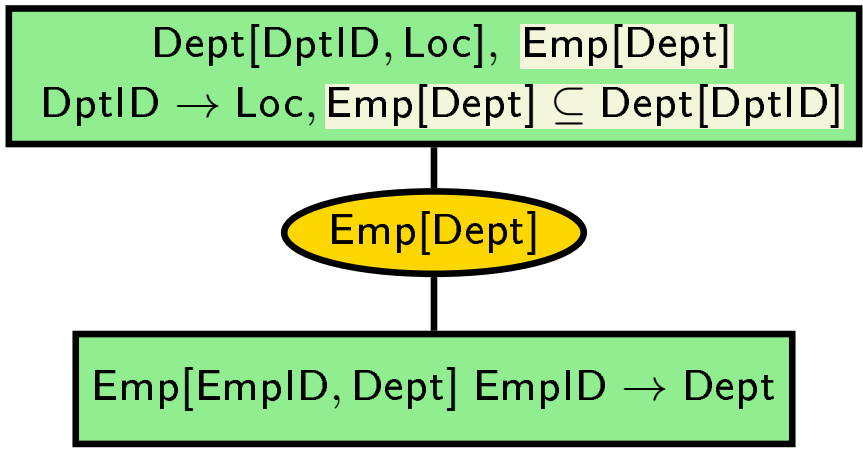

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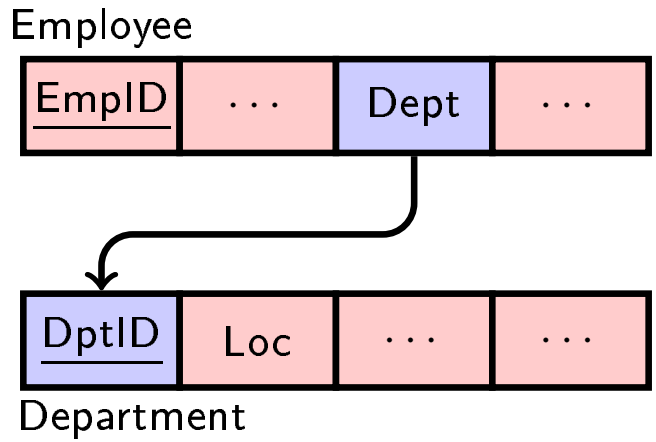

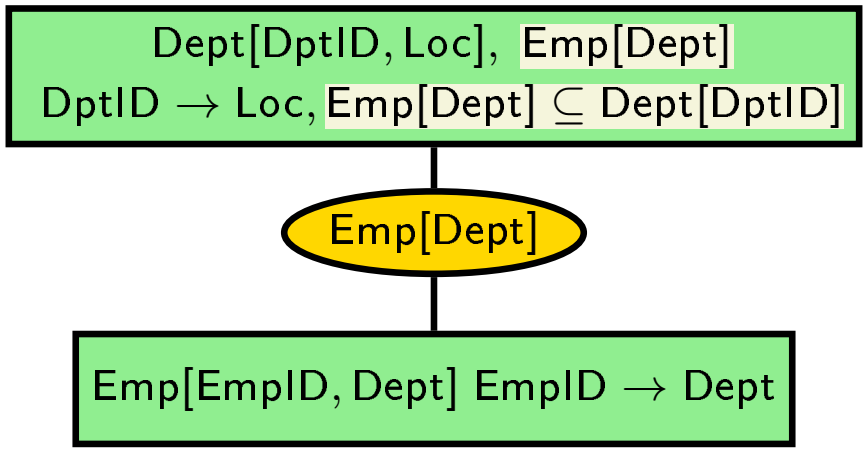

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