



Blocked In-Place Rectangular Transpose

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Main Idea

- Combine Cache Blocking with Point In-Place Transpose on a very tiny matrix
 - use of RB Format is the key idea
 - $CM \rightarrow RB \rightarrow RB^T \rightarrow CM$
 - SB is special case of RB
- Block In-Place Transpose is Very Fast relative to Point In-Place Transpose
- $CM \leftrightarrow RB$ uses fast vector In-Place Alg.

Summary or Overview

- A is M by N .
- $M = m * MB$ & $N = n * NB$
- $CM \rightarrow RB$ by vector IP transpose
- $RB \leftrightarrow RB^T$ by block IP transpose
 - use point IP transpose on m by n $A1 = SB A$
- $RB^T \rightarrow CM$ by vector IP transpose

Vector In-Place Xpose or $CM \leftrightarrow SB$

- Let A be M by NB with $M = m * MB$
- View A as m by NB A_1 with each $a_1(i,j)$ being a column vector of size NB
- Apply point IP transpose to A_1 to get A_2
- A_2 is m order NB SB 's concatenated
- Apply above subroutine $n = N / NB$ times

Where does the Speed Come From

- Data moved in blocks and vectors gives a 10 to 100 times performance gain
 - uses stride one processing; every line gets fully used when it enters L1 and streaming by algorithmic / automatic pre-fetching works
- SMP parallelism is easy to implement
 - disjoint cycle structure
 - long cycles can be broken into pieces

Other Matrix Layouts

- Can block transform (in-place) any permutation that can be described by a compact functional description
 - includes all common matrix data layouts
 - standard CM / RM rectangular arrays
 - standard CM / RM triangular arrays
 - standard packed format

An Example of CM to RB

- A is $M = 500$ by $NB = 4$
 - $M = m * MB$ with $MB = 100$ and $m = 5$
- A1 is $m = 5$ by $NB = 4$; each element of A1 is a vector of length $MB = 100$
- Both A and A1 are identical in storage and occupy $M * NB = 2000$ contiguous locations

Picture of the Previous slide

	0	500	1000	1500
	0	1	2	3
0	00	01	02	03
100	10	11	12	13
200	20	21	22	23
300	30	31	32	33
400	40	41	42	43

	0	400	800	1200	1600
	0	1	2	3	4
0	00	10	20	30	40
100	01	11	21	31	41
200	02	12	22	32	42
300	03	13	23	33	43

Previous slide shows CM to RB is vector in place transpose

- The left matrix $A1$ is a $m = 5$ by $NB = 4$ matrix whose elements are vectors of length $MB = 100$. This matrix is in standard CM format.
- The right matrix is the $NB = 4$ by $m = 5$ vector transpose of $A1$. It is in RB format consisting of $m = 5$ RB's of size $MB = 100$ by $NB = 4$.

Details of vector in place transpose

- 0 and $m*n - 1 = 10$ are singleton cycles
- 19 is prime and $\# d = 2; 1 \& 19$
- $q = m*n - 1 = 19$ is the mod value
- For problem 19, $\phi = 18 \& cl = 9$; leaders are 1, 2
- For problem 1, $\phi = 1 \& cl = 1$ at 19
- Further details follow

More Details continued

- cycle one = 1, 4, 16, 7, 9, 17, 11, 6, 5, 1
- cycle two = 2, 8, 13, 14, 18, 15, 3, 12, 10, 2
- cycle three = 19
- cycle four = 0
- These four cycles cover all of A1's twenty vectors of length $MB = 100$
- These four cycles cover all 2000 elements of A and transform A from CM format to RB format

A1 layout of A; see slide # 8

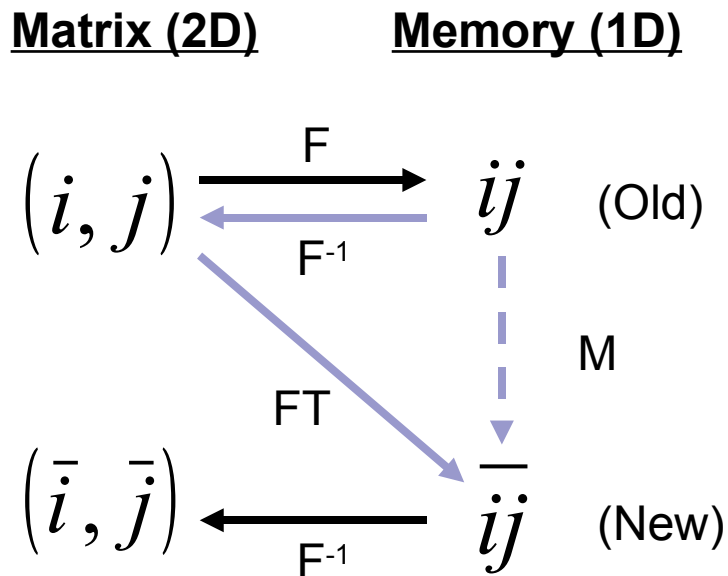
0 5 10 15
0 1 2 3

0	00	01	02	03
1	10	11	12	13
2	20	21	22	23
3	30	31	32	33
4	40	41	42	43

0 4 8 12 16
0 1 2 3 4

0	00	10	20	30	40
1	01	11	21	31	41
2	02	12	22	32	42
3	03	13	23	33	43

In Place Transpose Mapping



$$F(i, j) \mapsto i + m \cdot j$$

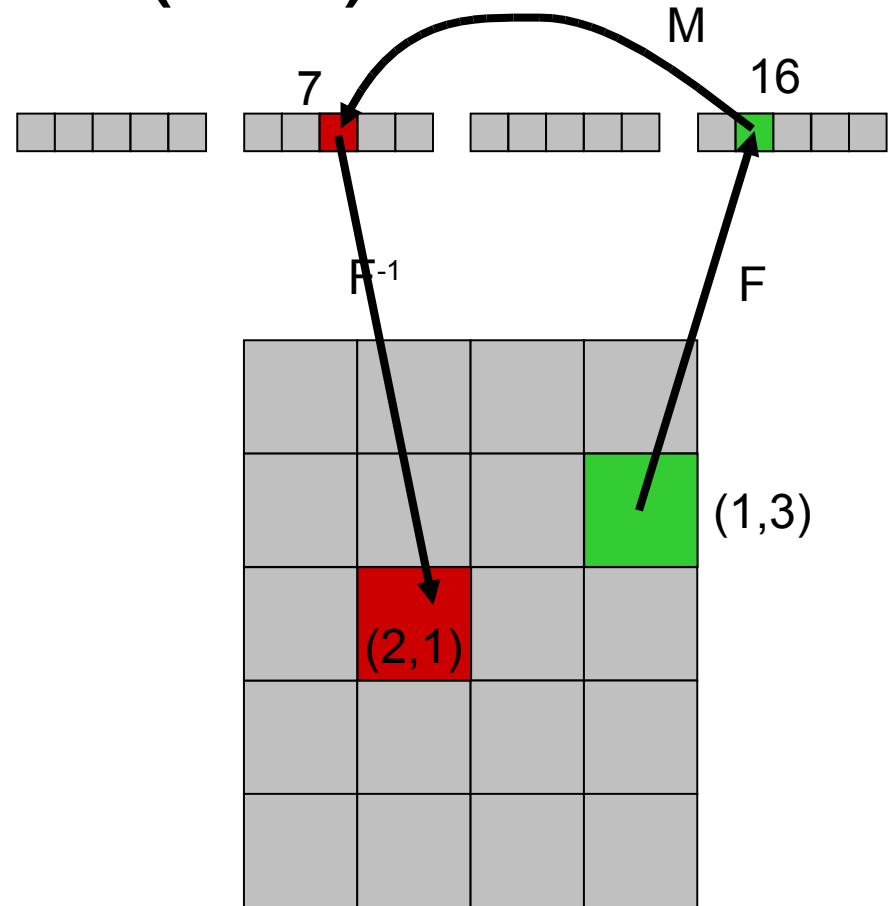
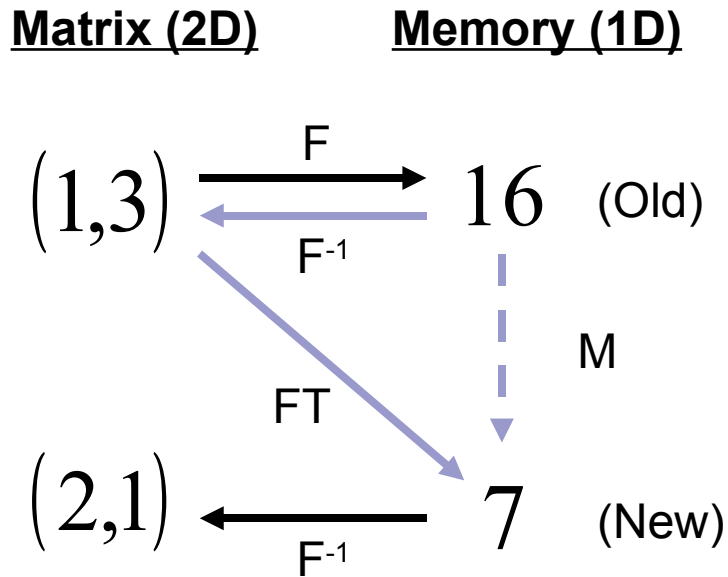
$$F^{-1}(ij) \mapsto \left(\text{mod}(ij, m), \left\lfloor \frac{ij}{m} \right\rfloor \right)$$

$$FT(i, j) \mapsto j + n \cdot i$$

$$M(ij) \mapsto \text{mod}(n \cdot ij, q)$$

$$M = FT \circ F^{-1}$$

Mapping (1,3) to (2,1)



Picture of the 4 cycles of slide # 11

0	8	8	5
0	7	6	2
0	3	7	5
6	1	2	4
1	4	3	0

A is 500 by 700 in CM order

- CM A has LDA = 500
- A has 7 column swaths: 500 by 100 each
- A1 is 5 by 100 matrix of vectors
- In-place transpose with $q = 499$
- repeat above 6 more times
- A is now in SB format of size 5 by 7

Details of CM to SB Vector

- 0 and $m*n - 1 = 499$ are singleton cycles
- 499 is prime and $\# d = 2; 1 \text{ \& } 499$
- $q = m*n - 1$ is the mod value
- for problem 499, $\phi = 498 \text{ \& } cl = 249$;
leaders are 1, 2
- for problem 1, $\phi = 1 \text{ \& } cl = 1$ at 499

Details of SB to SB^T

- $q = 5*7 - 1 = 34 = 2*17$
- $q = \text{sum over divisors of phi}$
 - $\#d = 4; 34, 17, 2, 1; \text{phi's} = 16, 16, 1, 1$
- $\#d$ problems gives cycles of length 16, 16, 1, 1 starting at 1, 2, 17, 34

Cycles for SB to SB^T

- Map $(ij) = \text{mod} (ij*n, q)$; $m = 5, n = 7,$
 $q = m*n - 1 = 34$
- cycle one : 1, 7, 15, 3, 21, 11, 9, 29, 33,
27, 19, 31, 13, 23, 25, 5, 1
- cycle two : 2, 14, 30, 6, 8, 22, 18, 24, 32,
20, 4, 28, 26, 12, 16, 10, 2
- cycles at 17, 34, and 0 are singletons

Picture of Map on slide # 19

0	15	15	2	9	14	2
0	3	5	14	4	12	11
0	1	13	0	5	9	8
3	4	12	6	13	11	8
10	6	1	10	7	7	0

Details of SB^T to CM

- $m = 100, n = 7, q = 699 = 3 \cdot 233$
- # d = 4; 699, 233, 3, 1; phi's 464, 232, 2, 1
- cl's are 166, 166, 1, 1
- leaders are 1, 2, 5, 10; 3, 9; 233, 466; 699

The 500 by 700 A as a point matrix

- $q = m \cdot n - 1 = 349,999 = 13^2 \cdot 19 \cdot 109$
- # d's = $3 \cdot 2 \cdot 2 = 12$:
- sum of $\phi(d) = q$
- twelve ϕ 's are 303264, 23328, 16848, 2808, 1944, 1296, 216, 156, 108, 18, 12, 1
- twelve cl 's are 468, 36, 156, 468, 18, 12, 36, 156, 6, 9, 12, 1
- ratio's give # of leaders: 648, 648, 108, 6, 108, 108, 6, 1, 18, 2, 1, 1: sum = 1655

500 by 700 A as point matrix

- hand-out has cycle of length 12 at $ij=247$
 - ij is $(247,0)$ element of A ; next element in cycle is $\text{mod}(247*700,q)$; $247 \mid q$ so we get cycle is $i \leftarrow \text{mod}(700*i,1417)$:
 - $247*(1, 700, 1135, 980, 172, 1372, 1091, 1354, 1244, 762, 608, 500, 1) \text{ mod } (q)$